

# 1 Syntax

$e ::= c \in \mathbb{Z}$ $\quad   x \in \text{Var}$ $\quad   e_1 + e_2$ $\quad   \lambda x. s$ $\quad   e_1(e_2)$	$e_e ::= \cdot +_1 e$ $\quad   \cdot +_2 \cdot$ $\quad   @_1(e_2)$ $\quad   @_2$ $\quad   @_3$	$s \in \text{stat} ::= \text{skip}$ $\quad   s_1; s_2$ $\quad   x := e$ $\quad   \text{if } (e > 0) s_1 s_2$ $\quad   \text{while } (e > 0) s$ $\quad   \text{return } e$	$s_e ::= x :=_1 \cdot$ $\quad   \cdot;_1 s_2$ $\quad   \text{if}_1 s_1 s_2$ $\quad   \text{while}_1 (e > 0) s$ $\quad   \text{while}_2 (e > 0) s$ $\quad   \text{return}_1 \cdot$
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# 2 Semantics

## 2.1 Expressions

$$\frac{\text{RED-CONST}(c)}{H_e, \ell_e, \ell_c, c \Downarrow H_e, \ell_e, c} \quad \frac{\text{RED-VAR-LOCAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, \ell_c[x]} \quad x \in \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-VAR-GLOBAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, E[x]} \quad x \in \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-VAR-UNDEF}(x)}{H_e, \ell_e, \ell_c, x \Downarrow \text{err}} \quad x \notin \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-ADD}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, \cdot +_1 e_2 \Downarrow r'}{H_e, \ell_e, \ell_c, e_1 + e_2 \Downarrow r'} \quad \frac{\text{RED-ADD-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad v_1, r, \cdot +_2 \cdot \Downarrow r'}{\ell_c, (H_e, \ell_e, v_1), \cdot +_1 e_2 \Downarrow r'}$$

$$\frac{\text{RED-ADD-2}}{v_1, (H_e, \ell_e, v_2), \cdot +_2 \cdot \Downarrow H_e, \ell_e, v_1 + v_2} \quad \frac{\text{RED-LAMBDA}(x, s)}{H_e, \ell_e, \ell_c, \lambda x. s \Downarrow H_e, \ell_e, (\ell_c, \lambda x. s)}$$

$$\frac{\text{RED-APP}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, @_1(e_2) \Downarrow r'}{H_e, \ell_e, \ell_c, e_1(e_2) \Downarrow r'} \quad \frac{\text{RED-APP-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad \ell'_c, x, s, r, @_2 \Downarrow r'}{\ell_c, (H_e, \ell_e, (\ell'_c, \lambda x. s)), @_1(e_2) \Downarrow r'}$$

$$\frac{\text{RED-APP-2}(s)}{H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c, s \Downarrow r \quad r, @_3 \Downarrow r'}{\ell_c, x, s, (H_e, \ell_e, v), @_2 \Downarrow r'} \quad \frac{\text{RED-APP-3-RET}}{ret(H_e, \ell_e, v), @_3 \Downarrow H_e, \ell_e, v}$$

$$\frac{\text{RED-APP-3-NO-RET}}{H_e, \ell_e, \ell_c, @_3 \Downarrow \text{err}}$$

## 2.2 Statements

$$\begin{array}{c}
\text{RED-SKIP} \\
\hline
H_e, \ell_e, \ell_c, \text{skip} \Downarrow H_e, \ell_e, \ell_c \\
\\
\text{RED-SEQ}(s_1, s_2) \\
\hline
H_e, \ell_e, \ell_c, s_1 \Downarrow r \quad r, \cdot;_1 s_2 \Downarrow r' \\
\hline
H_e, \ell_e, \ell_c, s_1; s_2 \Downarrow r' \\
\\
\text{RED-SEQ-1}(s_2) \\
\hline
H_e, \ell_e, \ell_c, s_2 \Downarrow r \\
\hline
H_e, \ell_e, \ell_c, \cdot;_1 s_2 \Downarrow r \\
\\
\text{RED-ASN}(x, e) \\
\hline
H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, x :=_1 \cdot \Downarrow r' \\
\hline
H_e, \ell_e, \ell_c, x := e \Downarrow r' \\
\\
\text{RED-ASN-1}(x) \\
\hline
\ell'_e = \text{fresh}(H_e) \quad E = H_e[\ell_e] \\
\hline
\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_e \leftarrow E[x \leftarrow v]], \ell'_e, \ell_c \quad x \notin \text{dom}(H_e[\ell_c]) \\
\\
\text{RED-ASN-1-LOCAL}(x) \\
\hline
\ell'_c = \text{fresh}(H_e) \quad C = H_e[\ell_c] \\
\hline
\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c \quad x \in \text{dom}(C) \\
\\
\text{RED-IF}(e, s_1, s_2) \\
\hline
H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{if}_1 s_1 s_2 \Downarrow r' \\
\hline
H_e, \ell_e, \ell_c, \text{if}(e > 0) s_1 s_2 \Downarrow r' \\
\\
\text{RED-IF-1-POS}(s_1, s_2) \\
\hline
H_e, \ell_e, \ell_c, s_1 \Downarrow r \\
\hline
\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r \quad v > 0 \\
\\
\text{RED-IF-1-NEG}(s_1, s_2) \\
\hline
H_e, \ell_e, \ell_c, s_2 \Downarrow r \\
\hline
\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r \quad v \leq 0 \\
\\
\text{RED-WHILE}(e, s) \\
\hline
H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{while}_1(e > 0) s \Downarrow r' \\
\hline
H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r' \\
\\
\text{RED-WHILE-1-NEG}(e, s) \\
\hline
\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow H_e, \ell_e, \ell_c \quad v \leq 0 \\
\\
\text{RED-WHILE-1-POS}(e, s) \\
\hline
H_e, \ell_e, \ell_c, s \Downarrow r \quad r, \text{while}_2(e > 0) s \Downarrow r' \\
\hline
\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow r' \quad v > 0 \\
\\
\text{RED-WHILE-2}(e, s) \\
\hline
H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r \\
\hline
H_e, \ell_e, \ell_c, \text{while}_2(e > 0) s \Downarrow r \\
\\
\text{RED-RETURN}(e) \\
\hline
H_e, \ell_e, \ell_c, e \Downarrow r \quad r, \text{return}_1 \cdot \Downarrow r' \\
\hline
H_e, \ell_e, \ell_c, \text{return } e \Downarrow r' \\
\\
\text{RED-RETURN-1} \\
\hline
(H_e, \ell_e, v), \text{return}_1 \cdot \Downarrow \text{ret}(H_e, \ell_e, v)
\end{array}$$

### 2.3 Aborting Rules

$$\begin{array}{c}
 \text{RED-ERROR-EXPR}(e) \\
 \hline
 \sigma, e \Downarrow \text{err} \quad \mathbf{abort} \sigma \wedge \mathbf{-intercept}_e \sigma
 \end{array}
 \qquad
 \begin{array}{c}
 \text{RED-ERROR-STAT}(s) \\
 \hline
 \sigma, s \Downarrow \text{err} \quad \mathbf{abort} \sigma
 \end{array}$$
  

$$\begin{array}{c}
 \sigma = C[\text{err}] \\
 \hline
 \mathbf{abort} \sigma
 \end{array}
 \qquad
 \frac{}{\mathbf{intercept}_{@_3} \text{ret}(H_e, \ell_e, v)}$$